# Identifying and Controlling Risky Contingencies of Transmission Systems 

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N-K criterion revisited

Previous work: Salmeron and Wood, Donde et al, Turitsyin, Hines

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- How about N - K, for K "larger"? Everybody knows that:
- It is too slow. A very difficult combinatorial problem.


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| ---: | ---: | ---: | ---: | :---: |
| $N$ | $K=2$ | $K=3$ | $K=4$ |  |
| 1000 | 499500 | 166167000 | 41417124750 |  |
| 4000 | 7998000 | 10658668000 | 10650673999000 |  |
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- Perhaps N - K does not necessarily capture all interesting events?


## Example: August 142003

U.S. - Canada report on blackout:
"Because it had been hot for several days in the Cleveland-Akron area, more air conditioners were running to overcome the persistent heat, and consuming relatively high levels of reactive power - further straining the area's limited reactive generation capabilities."
$\rightarrow$ A system-wide condition that impedes the system
$\rightarrow$ Not a cause, but a contributor
$\rightarrow$ Look for combined events ?

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- Perhaps N - K does not necessarily capture all interesting events?
- How can we deal with both types of problems?


## A continuous interdiction model

- A fictitious adversary is trying to interdict the transmission system.
- This adversary negatively alters the physical parameters of equipment, e.g. transmission lines, so as to impede transmission.
- The adversary has a budget available (both system-wide and per-line).
- On line $k m$, reactance $\boldsymbol{x}_{\boldsymbol{k} \boldsymbol{m}}$ increased to $\left(\mathbf{1}+\boldsymbol{\lambda}_{\boldsymbol{k m}}\right) \boldsymbol{x}_{\boldsymbol{k m}}$


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$-\sum_{k m} \boldsymbol{\lambda}_{\boldsymbol{k m}} \leq \boldsymbol{\Lambda}$ (global limit)


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- This adversary negatively alters the physical parameters of equipment, e.g. transmission lines, so as to impede transmission.
- The adversary has a budget available (both system-wide and per-line).
- Adversary maximizes the impact (e.g. voltage loss) over the available budget.
- A continuous, non-convex optimization problem with simple constraints.
No emumeration!


## A blast from the past: Bienstock and Verma, 2007

- DC approximation to power flows.
- Adversary increases reactances of lines.
- Limit on total percentage-increase of reactances, and on perline increase.
- Adversary maximizes the maximum line overload:

$$
\begin{aligned}
\max _{\boldsymbol{x}, \boldsymbol{\theta}} & \max _{k m}\left\{\frac{\left|\theta_{k}-\theta_{m}\right|}{u_{k m} x_{k m}}\right\} \\
\text { s.t. } & \boldsymbol{B}_{x} \theta=d \\
& \boldsymbol{x} \text { within budget }
\end{aligned}
$$

- Variables: reactances $x$, phase angles $\pi$
$-x_{k m}=$ reactance of $\boldsymbol{k m}, \boldsymbol{u}_{\boldsymbol{k m}}=$ limit of $\boldsymbol{k m}, \quad \boldsymbol{B}_{x}=$ bus susceptance matrix, $\boldsymbol{d}=$ net injections (given)
- Continuous, but non-smooth problem.

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& \sum_{k m}\left(\alpha_{k m}^{+}+\alpha_{k m}^{-}\right)=1, \quad \alpha^{+}, \alpha^{-} \geq 0 .
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- Continuous, smooth, nonconvex.


## Technical point

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Function to maximize: $\boldsymbol{F}(\boldsymbol{x}, \boldsymbol{\alpha}) \doteq \sum_{k m}\left(\alpha_{k m}^{+}-\alpha_{k m}^{-}\right) \frac{\left(\theta_{k}-\theta_{m}\right)}{u_{k m} x_{k m}}$

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- Fact: The gradient and the Hessian of $F(x, \alpha)$ can be efficiently computed
- Optimization problem solved using LOQO ( IPOPT an option)


## And what happens?

- Algorithm scales well (2007): CPU times of $\sim 1$ hour for studying systems with thousands of lines.


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- Optimal * attack concentrated on a handful of lines


## And what happens?

- Algorithm scales well (2007): CPU times of $\sim 1$ hour for studying systems with thousands of lines.
- Optimal * attack concentrated on a handful of lines
- Significant part of the budget expended on many lines, with visible impact

Table 6: Attack patterns

| single $=20$ |  |  |  |  |  |  |  | total $=\mathbf{6 0}$ | single $=\mathbf{1 0}$ |  | total $=\mathbf{3 0}$ | single $=\mathbf{1 0}$ | total $=\mathbf{4 0}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Range | Count | Range | Count | Range | Count |  |  |  |  |  |  |  |  |
| $[1,1]$ | 8 | $[1,1]$ | 1 | $[1,1]$ | 14 |  |  |  |  |  |  |  |  |
| $(1,2]$ | 72 | $(1,2]$ | 405 | $(1,2]$ | 970 |  |  |  |  |  |  |  |  |
| $(2,3]$ | 4 | $(2,9]$ | 0 | $(2,5]$ | 3 |  |  |  |  |  |  |  |  |
| $(5,6]$ | 1 | $(9,10]$ | 3 | $(5,6]$ | 0 |  |  |  |  |  |  |  |  |
| $(6,7]$ | 1 |  |  | $(6,7]$ | 1 |  |  |  |  |  |  |  |  |
| $(7,8]$ | 4 |  |  | $(7,9]$ | 0 |  |  |  |  |  |  |  |  |
| $(8,20]$ | 0 |  |  |  |  |  |  |  |  |  |  |  |  |

"single" $=$ max multiplicative increase of a line's reactance
"total" = max total multiplicative increase of line reactances

## Today: the AC power flows setting

As before, adversary increases impedances, subject to budgets
Adversary wants to maximize:

- Phase angle differences across ends of a lines
- Voltage deviations (loss)

Alternative version:

- There is a recourse action: shed load so as to maintain feasibility of all power flow constraints (limits)
- Adversary wants to maximize the amount of lost load

Generically:

$$
\begin{aligned}
\max & \mathcal{F}(x) \\
\text { s.t. } & x \in \mathcal{B}
\end{aligned}
$$

- $\boldsymbol{x}=$ impedances, $\boldsymbol{\mathcal { B }}=$ budget constraints
- $\mathcal{F}(x)=$ meausure of phase angle differences, voltage loss, load loss
- Challenge 1: $\mathcal{F}(x)$ is implicitly defined

Basic methodology: Frank-Wolfe

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$$
\begin{array}{rll}
\boldsymbol{y}^{k} \text { solves } & \max & {\left[\nabla \mathcal{F}\left(x^{k}\right)\right]^{T} y} \\
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\end{array}
$$

Final step is a line search: $x^{k+1}=x^{k}+\alpha y^{k}$, where $\mathbf{0} \leq \boldsymbol{\alpha} \leq \mathbf{1}$ is the stepsize.

## Line searches






## Challenge 2



- Recall: $\mathcal{F}(x)$ measures e.g. the largest phase angle difference using reactances $\boldsymbol{x}$
- Q: exactly how do we get $\nabla \mathcal{F}(x)$ ?


## Challenge 2

```
\(\max \mathcal{F}(x)\)
    s.t. \(\quad x \in \mathcal{B}\) (within budget)
```



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- But $\nabla \mathcal{F}(x)$ is a vector with an entry for each line of the transmission system - it is a big vector
- "Solution": Estimate $\nabla \mathcal{F}(x)$ in parallel over several cores
- Alternative: only estimate some of the components of $\nabla \mathcal{F}(x)$ :
- Random subset of small size
- Most promising subset


## Challenge 3



- $\mathcal{F}(\boldsymbol{x})$ measures e.g. the sum of voltage losses with reactances $\boldsymbol{x}$
- And we estimate $\nabla \mathcal{F}(x)$ using finite differences
- Q: How do we compute $\mathcal{F}(\boldsymbol{x})$, for given reactances $\boldsymbol{x}$ ?


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- Challenge! PF often does not converge for interesting $x$
- solution: solution OPF-like problem:
minimize sum of square of all violations (load mismatch, line limits, etc)
- solution? violations still observed
- solution? Add to definition of $\mathcal{F}(x)$ sum of weighted square violations $\rightarrow$ Currently using IPOPT within Matpower (fastest for our purposes)
$\rightarrow$ Infeasible cases verified using SDP relaxation


## Example: phase angle attack on Polish grid (from Matpower)

```
1 obj=2620.72 step=1.00 [ 263 8.00; 300 8.00; 728 8.00; ]
2 obj=2641.52 step=1.00 [ 305 8.00; 306 8.00; 309 8.00; ]
3 obj=2649.34 step=1.00 [ 168 8.00; 263 8.00; 321 8.00; ]
5 obj=2765.47 step=0.50[51 4.00; 261 4.00; 263 4.00; 300 4.00; 321 4.00; 322 4.00;]
13 obj=2944.01 step=0.12 [ 305 2.60; 168 2.32; 322 2.17; 169 1.90; 321 1.85; 263 1.57; 309
1.50; 32 1.15; 51 1.08; 261 1.08; 170 1.00; 171 1.00; 306 0.85; 39 0.75; 281 0.75; 166 0.57;
310 0.57; 8 0.43; 264 0.43; 300 0.42; ]
20 obj=2950.54 step=0.03[169 2.53; 305 2.38; 168 1.88; 322 1.77; 321 1.76; 309 1.74; 166
1.44; 170 1.28; 263 1.28; 261 1.14; 32 0.93; 51 0.88; 171 0.81; 306 0.69; 39 0.61; 281 0.61;
264 0.59; 260 0.51; 310 0.46; 8 0.35; 300 0.34; ]
27 obj=2958.08 step=0.00 [ 169 2.80; 305 2.53; 321 2.00; 309 1.97; 168 1.63; 263 1.58; 322
1.53; 166 1.38; 261 1.11; 170 1.11; 32 0.81; 51 0.76; 264 0.76; 281 0.75; 171 0.71; 306 0.60;
39 0.53; 260 0.44; 310 0.40; 8 0.30; 300 0.30; ]
```


## Example: phase angle attack on 118-bus

Three top-attacked lines in red:


Fact: phase angle attack cannot be isolated to a few lines

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Experiment on 118-bus case:
(1) Take line most heavily interdicted: line $\mathbf{3 8}$
(2) Let the reactance of this line increase to infinity
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From $\approx 10$ to $\approx 40$.


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From $\approx 10$ to $\approx 40$.


## Voltage attack on 118-bus

"Triple the reactance of at most three lines"
Voltage Magnitude


Changes in Voltage Magnitude



## Voltage attack on 2383-bus Polish

"Double the reactance of at most three lines"


Changes in Voltage Magnitude

$\rightarrow$ Primarily 4 lines interdicted

