Identifying and Controlling Risky Contingencies of Transmission Systems

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FERC Software conference, 2015

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Previous work: Salmeron and Wood, Donde et al, Turitsyin, Hines

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1000	499500	166167000	41417124750
4000	7998000	10658668000	10650673999000
8000	31996000	85301336000	170538695998000
10000	49995000	166616670000	416416712497500

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- Perhaps N K does not necessarily capture all interesting events?

Example: August 14 2003

U.S. - Canada report on blackout:

"Because it had been hot for several days in the Cleveland-Akron area, more air conditioners were running to overcome the persistent heat, and consuming relatively high levels of reactive power – further straining the area's limited reactive generation capabilities."

- → A **system-wide** condition that impedes the system
- → Not a cause, but a contributor
- \rightarrow Look for combined events?

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 (T. Boston) during Hurricane Sandy, N 142 was observed.
- Perhaps N K does not necessarily capture all interesting events?
- How can we deal with both types of problems?

- A fictitious adversary is trying to interdict the transmission system.
- This adversary negatively alters the physical parameters of equipment, e.g. transmission lines, so as to impede transmission.
- The adversary has a budget available (both system-wide and per-line).
 - On line km, reactance x_{km} increased to $(1 + \lambda_{km})x_{km}$

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 - $-\sum_{km} \lambda_{km} \leq \Lambda \text{ (global limit)}$

- A fictitious adversary is trying to interdict the transmission system.
- This adversary negatively alters the physical parameters of equipment, e.g. transmission lines, so as to impede transmission.
- The adversary has a budget available (both system-wide and per-line).
- Adversary maximizes the impact (e.g. voltage loss) over the available budget.
- A continuous, non-convex optimization problem with **sim- ple** constraints.

No emumeration!

A blast from the past: Bienstock and Verma, 2007

- DC approximation to power flows.
- Adversary increases reactances of lines.
- **Limit** on total percentage-increase of reactances, and on perline increase.
- Adversary maximizes the maximum line overload:

$$\max_{\boldsymbol{x},\boldsymbol{\theta}} \max_{km} \left\{ \frac{|\theta_k - \theta_m|}{u_{km} \boldsymbol{x_{km}}} \right\}$$
s.t.
$$\boldsymbol{B_x}\theta = d$$

$$\boldsymbol{x} \text{ within budget}$$

- Variables: reactances \boldsymbol{x} , phase angles $\boldsymbol{\pi}$
- $-\mathbf{x}_{km}$ = reactance of km, \mathbf{u}_{km} = limit of km, \mathbf{B}_{x} = bus susceptance matrix, \mathbf{d} = net injections (given)
- Continuous, but non-smooth problem.

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$$\sum_{km} (\alpha_{km}^{+} + \alpha_{km}^{-}) = 1, \quad \alpha^{+}, \alpha^{-} \geq 0.$$

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• Continuous, smooth, **nonconvex**.

Technical point

$$\max_{\boldsymbol{x},\boldsymbol{\theta},\boldsymbol{\alpha}} \sum_{km} (\alpha_{km}^{+} - \alpha_{km}^{-}) \frac{(\theta_{k} - \theta_{m})}{u_{km}} \frac{\boldsymbol{x}_{km}}{\boldsymbol{x}_{km}}$$
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Function to maximize:
$$F(x, \alpha) \doteq \sum_{km} (\alpha_{km}^+ - \alpha_{km}^-) \frac{(\theta_k - \theta_m)}{u_{km} x_{km}}$$

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Function to maximize: $F(x, \alpha) \doteq \sum_{km} (\alpha_{km}^+ - \alpha_{km}^-) \frac{(\theta_k - \theta_m)}{u_{km} x_{km}}$

- Fact: The gradient and the Hessian of $F(x, \alpha)$ can be efficiently computed
- Optimization problem solved using **LOQO** (**IPOPT** an option)

And what happens?

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- Algorithm scales well (2007): CPU times of ~ 1 hour for studying systems with thousands of lines.
- Optimal * attack concentrated on a handful of lines
- Significant part of the budget expended on many lines, with visible impact

Table 6: Attack patterns

single = 20	total = 60	single = 10	total = 30	single = 10	total = 40
Range	Count	Range	Count	Range	Count
[1, 1]	8	[1, 1]	1	[1, 1]	14
(1,2]	72	[1, 2]	405	(1, 2]	970
(2,3]	4	(2, 9]	0	(2, 5]	3
(5,6]	1	(9, 10]	3	(5, 6]	0
(6,7]	1			(6, 7]	1
(7,8]	4			(7, 9]	0
(8, 20]	0			(9, 10]	2

[&]quot;single" = max multiplicative increase of a line's reactance

[&]quot;total" = max total multiplicative increase of line reactances

Today: the AC power flows setting

As before, adversary increases impedances, subject to budgets

Adversary wants to **maximize:**

- Phase angle differences across ends of a lines
- Voltage deviations (loss)

Alternative version:

- There is a **recourse** action: shed load so as to maintain feasibility of all power flow constraints (limits)
- Adversary wants to maximize the amount of lost load

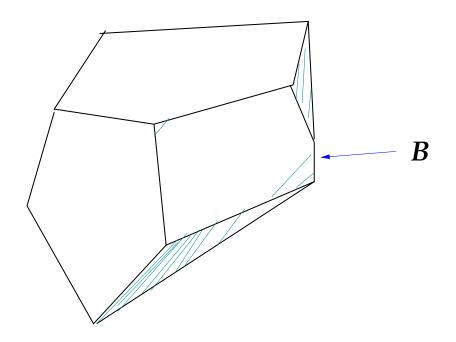
Generically:

$$\max_{\text{s.t.}} \quad \frac{\mathcal{F}(x)}{x \in \mathcal{B}}$$

- \mathbf{x} = impedances, $\mathbf{\mathcal{B}}$ = budget constraints
- $\mathcal{F}(x)$ = measure of phase angle differences, voltage loss, load loss
- Challenge 1: $\mathcal{F}(x)$ is implicitly defined

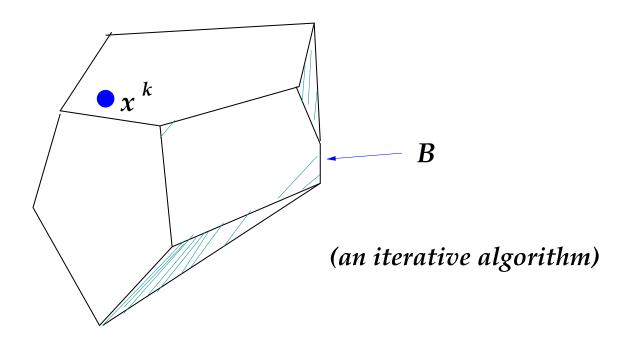
Basic methodology: Frank-Wolfe

 $egin{array}{ll} \max & \mathcal{F}(x) \\ \mathrm{s.t.} & x \in \mathcal{B} & \mathrm{(within budget)} \end{array}$



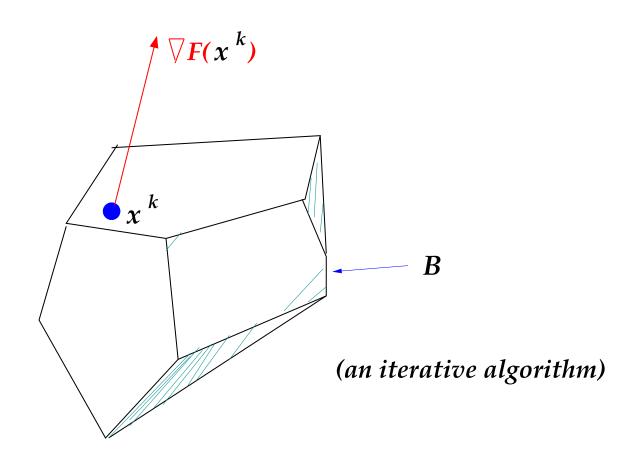
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 $\max \quad \frac{\mathcal{F}(x)}{s.t.}$ s.t. $x \in \mathcal{B}$ (within budget)



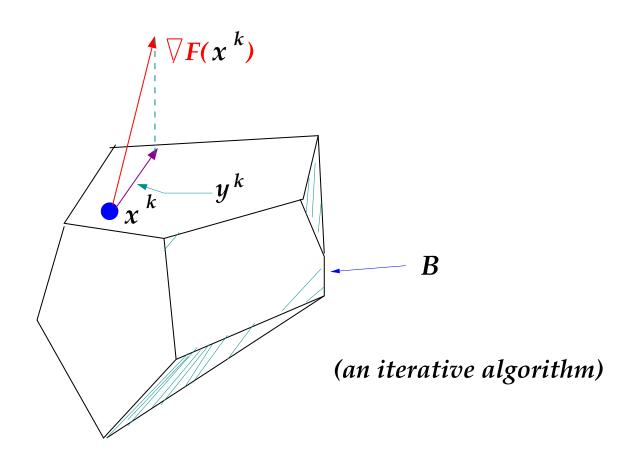
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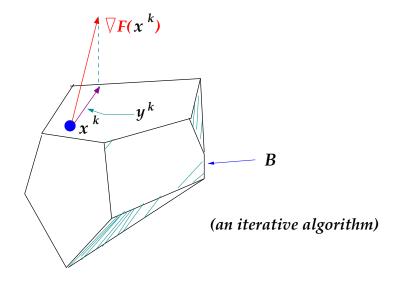
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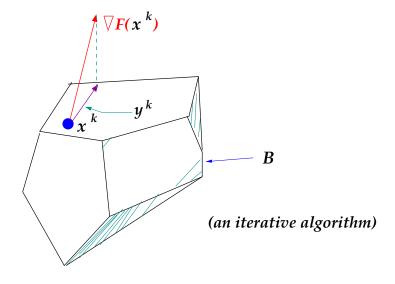
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$$m{y^k}$$
 solves $m{\max} [m{
abla} m{\mathcal{F}}(m{x^k})]^Tm{y}$ s.t. $m{x^k} + m{y} \in \mathcal{B}$ (within budget)

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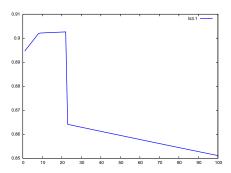
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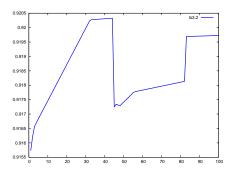


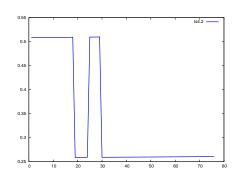
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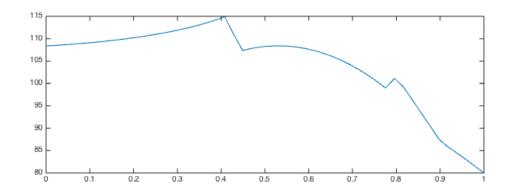
Final step is a line search: $x^{k+1} = x^k + \alpha y^k$, where $0 \le \alpha \le 1$ is the stepsize.

Line searches

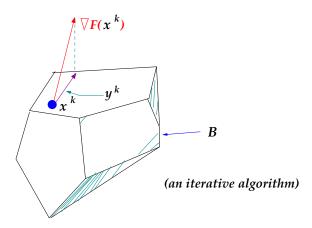






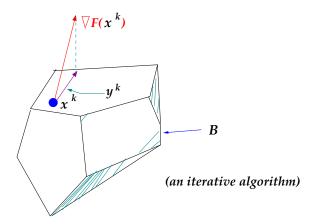


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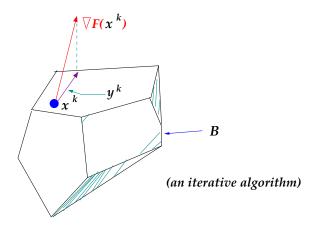


- ullet Recall: $m{\mathcal{F}}(m{x})$ measures e.g. the largest phase angle difference using reactances $m{x}$
- Q: exactly how do we get $\nabla \mathcal{F}(x)$?

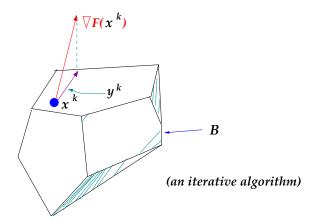
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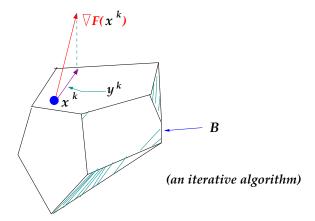
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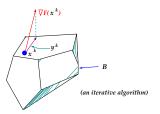
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- "Solution": Estimate $\nabla \mathcal{F}(x)$ in parallel over several cores

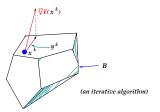


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- But $\nabla \mathcal{F}(x)$ is a vector with an entry for each line of the transmission system it is a **big** vector
- "Solution": Estimate $\nabla \mathcal{F}(x)$ in parallel over several cores
- Alternative: only estimate some of the components of $\nabla \mathcal{F}(x)$:
 - Random subset of small size
 - Most promising subset



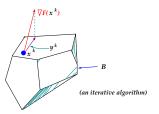
- $\bullet \mathcal{F}(x)$ measures e.g. the sum of voltage losses with reactances x
- \bullet And we estimate $\nabla \mathcal{F}(x)$ using finite differences
- Q: How do we compute $\mathcal{F}(x)$, for given reactances x?

 $\max \quad \frac{\mathcal{F}(x)}{\text{s.t.}} \quad \text{s.t.} \quad \frac{x}{\mathbf{e}} \in \mathcal{B} \quad \text{(within budget)}$



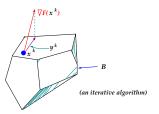
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- A: Ideally, a PF (load flow) calculation

 $\max \quad \mathcal{F}(x)$ s.t. $x \in \mathcal{B}$ (within budget)



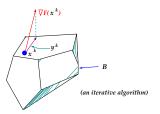
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- Challenge! PF often does not converge for interesting x

 $\max_{\mathbf{x}} \quad \frac{\mathcal{F}(\mathbf{x})}{\mathbf{x}}$ s.t. $\mathbf{x} \in \mathcal{B} \quad \text{(within budget)}$



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- solution: solution OPF-like problem: minimize sum of square of all violations (load mismatch, line limits, etc)

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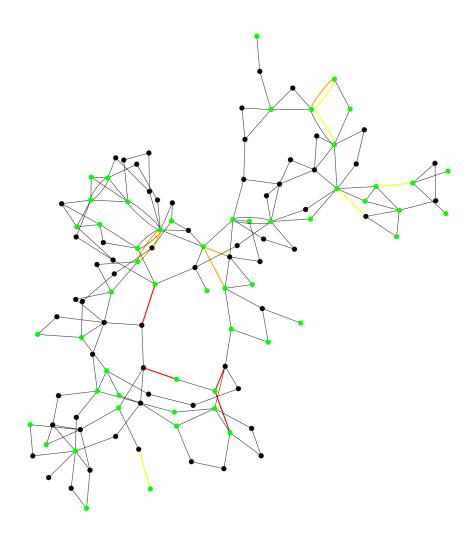
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- solution: solution OPF-like problem: minimize sum of square of all violations (load mismatch, line limits, etc)
- solution? violations still observed
- solution? Add to definition of $\mathcal{F}(x)$ sum of weighted square violations
- → Currently using **IPOPT** within Matpower (fastest for **our** purposes)
- → Infeasible cases verified using SDP relaxation

Example: phase angle attack on Polish grid (from Matpower)

```
1 obj=2620.72 step=1.00 [ 263 8.00; 300 8.00; 728 8.00; ]
2 obj=2641.52 step=1.00 [ 305 8.00; 306 8.00; 309 8.00; ]
3 obj=2649.34 step=1.00 [ 168 8.00; 263 8.00; 321 8.00; ]
5 obj=2765.47 step=0.50 [ 51 4.00; 261 4.00; 263 4.00; 300 4.00; 321 4.00; 322 4.00; ]
13 obj=2944.01 step=0.12 [ 305 2.60; 168 2.32; 322 2.17; 169 1.90; 321 1.85; 263 1.57; 309
1.50; 32 1.15; 51 1.08; 261 1.08; 170 1.00; 171 1.00; 306 0.85; 39 0.75; 281 0.75; 166 0.57;
310 0.57; 8 0.43; 264 0.43; 300 0.42;
20 obj=2950.54 step=0.03 169 2.53; 305 2.38; 168 1.88; 322 1.77; 321 1.76; 309 1.74; 166
1.44; 170 1.28; 263 1.28; 261 1.14; 32 0.93; 51 0.88; 171 0.81; 306 0.69; 39 0.61; 281 0.61;
264 0.59; 260 0.51; 310 0.46; 8 0.35; 300 0.34;
27 obj=2958.08 step=0.00 [ 169 2.80; 305 2.53; 321 2.00; 309 1.97; 168 1.63; 263 1.58; 322
1.53; 166 1.38; 261 1.11; 170 1.11; 32 0.81; 51 0.76; 264 0.76; 281 0.75; 171 0.71; 306 0.60;
39 0.53; 260 0.44; 310 0.40; 8 0.30; 300 0.30;
```

Example: phase angle attack on 118-bus

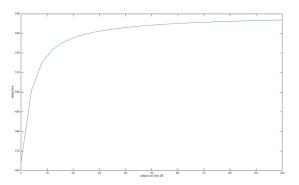
Three top-attacked lines in red:



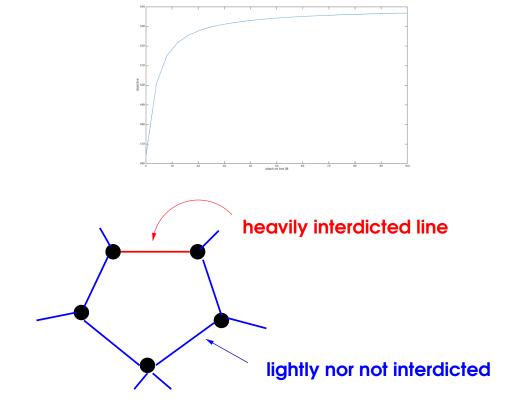
- (1) Take line most heavily interdicted: line **38**
- (2) Let the reactance of this line increase to infinity
- (3) What happens? Phase angle difference $\rightarrow \pi/2$?

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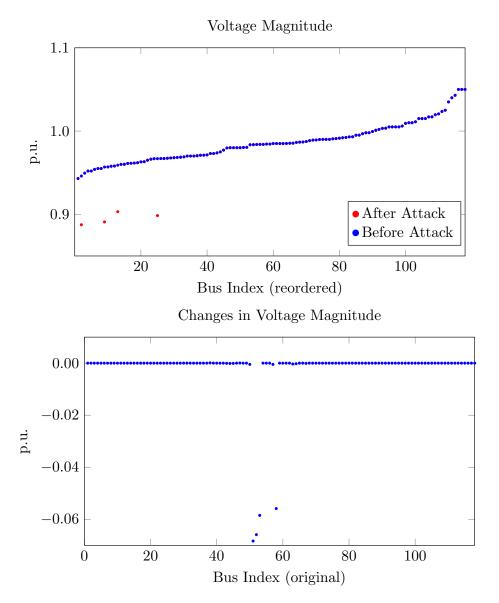


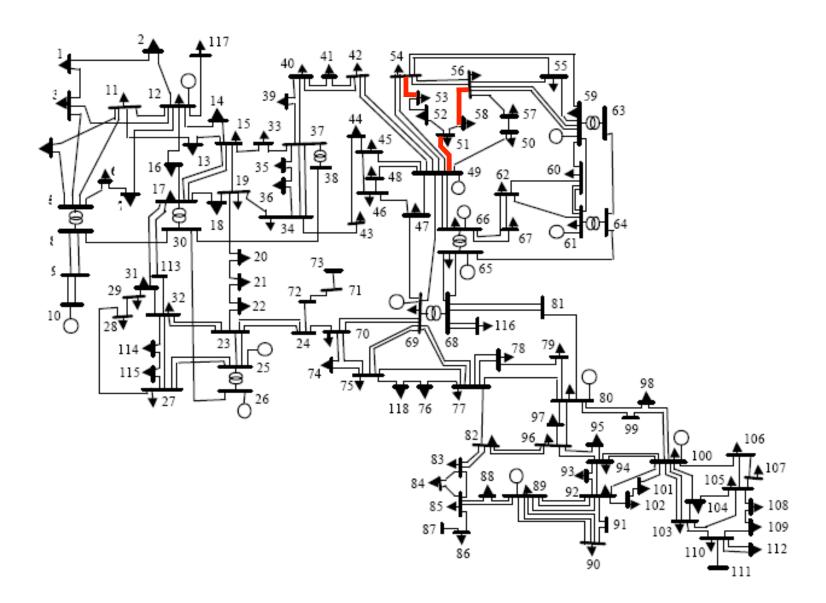
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Voltage attack on 118-bus

"Triple the reactance of at most three lines"





Voltage attack on 2383-bus Polish

"Double the reactance of at most three lines"

