Decentralized Robust Optimization Algorithms for Tie-Line Scheduling of Multi-Area Grid with Variable Wind Energy

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Outline

Background and Motivation

Robust Optimization Formulations and Properties

Distributed Computing Methods through Alternating Direction Method of Multipliers (ADMM)

□Numerical Study and Demonstrations

Conclusions

Part I: Background and Motivation

- Multi-area power system
 - Physically separated regions (China)
 - Different electricity market
- Intermittent renewable energy
- Random contingencies
- Aligned LMPs for boundary nodes



Challenges in Multi-Area Power Grids

- Systems are cooperative but operated independently
 - o Privacy
 - Commercial information
- Protocols between systems to support cooperation:
 - Coordinated Transaction Scheduling: NYISO and PJM
 - o Interchange Optimization: PJM and MISO
 - o Inter-Regional Interchange Scheduling: ISONE and NYISO
- Interfaces between multiple systems: tie-lines
 - Power flow decisions (day-ahead and intraday)
 - Unit commitment decisions (day-ahead)

Decision Models for Multi-Area Power Grids

Lagrangian relaxation

- Conejo and Aguado (1998), Multi-area coordinated decentralized DC optimal power flow
- Aguado, Quintana, and Conejo (1999), Optimal power flows of interconnected power systems
- Augmented Lagrangian decomposition
 - Kim and Baldick (1997)
 - Ahmadi-Khatir, Conejo, and Cherkaoui (2014)

Alternative Direction Method of Multiplers (ADMM)

- Kim and R. Baldick (2000)
- Zheng, Wu, Zhang, Sun, and Liu (2015)

Uncertainty Consideration in Power Grids

- Scenario based probabilistic models
 - Stochastic programming (SP) and chance constrained formulation
 - Explicit large-scale models but many existing computing methods: e.g., Ahmadi-Khatir, Conejo, and Cherkaoui (2013, 2014):
 - SP+ augmented Lagrangian relaxation
 - Issue: Inaccurate prediction of scenarios and probabilities -> infeasible solutions
- Robust optimization models
 - Uncertainty set based compact formulation
 - Produce highly feasible solutions by considering all possibilities inside uncertainty set
 - Two types of algorithms: Benders decomposition and *column-and-constraint generation*
 - Our aim: integration of decentralized computing scheme + robust optimization

Part II: Robust Optimization Formulations and Properties

Multi-Area Robust Tie-Line Scheduling

Two-stage decision making framework

$$\min_{\boldsymbol{\xi}^{f} \in \Omega^{f}} \sum_{a \in \mathcal{A}} \left\{ \max_{\tilde{\mathbf{P}}_{a}^{w} \in \mathbf{U}_{a}^{w}} \left[\min_{\boldsymbol{\xi}_{a}^{s} \in \Omega_{a}^{s}(\boldsymbol{\xi}_{a}^{f}, \tilde{\mathbf{P}}_{a}^{w})} C_{a}^{ED}(\boldsymbol{\xi}_{a}^{s}, \tilde{\mathbf{P}}_{a}^{w}) \right] \right\}$$

- Tie-line interchanges: the first-stage decisions before the availability of wind energy is known
 - Phase angles at boundary buses:

$$\begin{split} \boldsymbol{\xi}^{f} &= \{\boldsymbol{\xi}^{f}_{a}, \forall a \in \mathbf{A}\}, \boldsymbol{\xi}^{f}_{a} = \left\{\boldsymbol{\delta}_{a,i,t}, \forall g \in \mathbf{G}_{a}, i \in \mathbf{N} \left| {}^{BB}_{a} \cup \mathbf{N} \left| {}^{'BB}_{a}, t \in \mathbf{T} \right. \right\} \\ \Omega^{f} &= \left\{ \boldsymbol{\xi}^{f} \left| \boldsymbol{\delta}_{\phi(i),i,t} = \boldsymbol{\delta}_{\phi(j),i,t}, \right. \left| \boldsymbol{\delta}_{\phi(i),j,t} = \boldsymbol{\delta}_{\phi(j),j,t}, \left| \boldsymbol{\delta}_{\phi(i),i,t} - \boldsymbol{\delta}_{\phi(i),j,t} \right| \right| / X_{i,j} \leq \overline{F}_{i,j}, \forall (i,j) \in \mathbf{E}^{tie}, i > j, t \in \mathbf{T}, \\ \boldsymbol{\delta}_{a',ref,t} = 0, \forall a' \in \Phi(ref), t \in \mathbf{T} \right\} \end{split}$$

Economic dispatch: the second-stage decision after wind energy is revealed.
Continuous model

Inter-regional constraints

Region-coupling constraints - perceived phase angles are same at two ends of a tie-line



$$\delta_{a,i,t} = \delta_{b,i,t}, \, \delta_{a,j,t} = \delta_{b,j,t}, \, \forall t \in \mathbf{T}$$

Economic Dispatch of Each Area

- Dispatch of conventional generation units, wind farms, and phase angles of internal buses
- $\Omega_a^s(\boldsymbol{\xi}_a^f, \tilde{\mathbf{P}}_a^w) = \left\{\boldsymbol{\xi}_a^s\right\}$ $\sum_{i \in \Psi^{IB}(i)} (\theta_{i,t} - \theta_{j,t}) / X_{i,j} + \sum_{i \in \Psi^{BB}(i)} (\theta_{i,t} - \delta_{a,j,t}) / X_{i,j}$ $= \sum_{g \in \Psi_{a}^{G}(i)} p_{g,t}^{G} + \sum_{k \in \Psi_{a}^{W}(i)} p_{k,t}^{w} - \sum_{d \in \Psi^{D}(i)} p_{d,t}^{D}, \forall i \in \mathbb{N}_{a}^{B}, t,$ $\sum_{i \in \Psi^{IB}(i)} (\delta_{a,i,t} - \theta_{j,t}) / X_{i,j} + \sum_{i \in \Psi^{BB}(i)} (\delta_{a,i,t} - \delta_{a,j,t}) / X_{i,j}$ $=\sum_{g\in\Psi^G_+(i)}p^G_{g,t}+\sum_{k\in\Psi^W_+(i)}p^w_{k,t}-\sum_{d\in\Psi^D_+(i)}p^D_{d,t}, \forall i \in \mathbb{N}^{BB}_{a}, t,$ $-\overline{F}_{i,j} \leq (\theta_{i,t} - \theta_{j,t}) / X_{i,j} \leq \overline{F}_{i,j}, \forall i \in \mathbb{N}_{a}^{B}, j \in \Psi_{a}^{B}(i), j > i,$ $-\overline{F}_{i,i} \leq \left(\delta_{a,i,t} - \theta_{i,t}\right) / X_{i,j} \leq \overline{F}_{i,i}, \forall i \in \mathbb{N}_{a}^{BB}, j \in \Psi_{a}^{IB}(i),$

$$p_{g,t}^{G} + r_{g,t}^{G+} \leq P_{g}^{G}, \forall g \in \mathcal{G}_{a}, t \in \mathcal{T} ,$$

$$r_{g,t}^{G-} \ge p_{g,t}^G - \underline{P}_g, \forall g \in \mathbf{G}_a, t \in \mathbf{T}$$
,

$$\sum_{g \in \mathcal{G}_a} r_{g,t}^{G_+} \ge SR_{a,t}^+, \sum_{g \in \mathcal{G}_a} r_{g,t}^{G_-} \ge SR_{a,t}^-, \forall t \in \mathcal{T} ,$$

$$0 \leq r_{g,t}^{G_{+}} \leq RU_{g}^{G}, 0 \leq r_{g,t}^{G_{-}} \leq RD_{g}^{G}, \forall g \in G_{a}, t \in T,$$
$$0 \leq p_{k,t}^{w} \leq \tilde{P}_{k,t}^{w}, \forall k \in W_{a}, t \in T$$

Observation:

If the uncertainty sets U^w_a are polyhedra, the robust multi-area tie-line schedule problem is a convex optimization (an extremely large-scale linear program)

Idea: enumerating all extreme points of U^w_a to construct the equivalent formulation, which is an LP

Multi-Area Generation Unit and Tie-line Scheduling

Two-stage decision making framework

$$\min_{\boldsymbol{\xi}^{f} \in \Omega^{f}} \sum_{a \in \mathbf{A}} \left\{ C_{a}^{UD} \left(\boldsymbol{\xi}_{a}^{f}\right) + \max_{\tilde{\mathbf{P}}_{a}^{W} \in \mathbf{U}_{a}^{W}} \left[\min_{\boldsymbol{\xi}_{a}^{s} \in \Omega_{a}^{s} \left(\boldsymbol{\xi}_{a}^{f}, \tilde{\mathbf{P}}_{a}^{W}\right)} C_{a}^{ED} \left(\boldsymbol{\xi}_{a}^{s}, \tilde{\mathbf{P}}_{a}^{W}\right) \right] \right\}$$

• First-stage decisions: unit commitments and tie-line interchanges

$$\begin{split} \boldsymbol{\xi}^{f} &= \{\boldsymbol{\xi}_{a}^{f}, \forall a \in \mathbf{A}\}, \boldsymbol{\xi}_{a}^{f} = \{\boldsymbol{u}_{g,t}^{G}, \boldsymbol{x}_{g,t}^{G}, \boldsymbol{y}_{g,t}^{G}, \boldsymbol{\delta}_{a,i,t}, \forall g \in \mathbf{G}_{a}, i \in \mathbf{N} \left| \overset{BB}{a} \cup \mathbf{N} \right|_{a}^{BB}, t \in \mathbf{T} \} \\ \Omega^{f} &= \{\boldsymbol{\xi}^{f} \left| \forall a \in \mathbf{A}, g \in \mathbf{G}_{a}, t \in \mathbf{T}, \boldsymbol{\delta}_{\phi(i),i,t} = \boldsymbol{\delta}_{\phi(j),i,t}, \left| \boldsymbol{\delta}_{\phi(i),j,t} \right| \left| \boldsymbol{\delta}_{\phi(i),i,t} - \boldsymbol{\delta}_{\phi(i),j,t} \right| \right| \left| \boldsymbol{X}_{i,j} \leq \overline{F}_{i,j}, \forall (i,j) \in \mathbf{E}^{tie}, i > j, \\ \boldsymbol{\delta}_{a',ref,t} = 0, \forall a' \in \Phi(ref), t \in \mathbf{T}, \\ \boldsymbol{u}_{g,t}^{G} - \boldsymbol{u}_{g,t-1}^{G} = \boldsymbol{x}_{g,t}^{G} - \boldsymbol{y}_{g,t}^{G}, \sum_{\tau = \max\left\{1, t - MU_{g}^{G} + 1\right\}}^{t} \boldsymbol{x}_{g,\tau}^{G} \leq \boldsymbol{u}_{g,t}^{G}, \sum_{\tau = \max\left\{1, t - MU_{g}^{G} + 1\right\}}^{t} \boldsymbol{y}_{g,\tau}^{G} \leq 1 - \boldsymbol{u}_{g,t}^{G}, \forall g, t \\ \boldsymbol{u}_{g,t}^{G} \in \{0,1\}, 0 \leq \boldsymbol{x}_{g,t}^{G} \leq 1, 0 \leq \boldsymbol{y}_{g,t}^{G} \leq 1 \end{split}$$

Multi-Area Generation Unit and Tie-line Scheduling (Cont'd)

 Second-stage decisions: economic dispatch after the available wind power is revealed and unit status are determined

Observation:

 Due to binary variables for unit scheduling, the robust formulation is equivalent to a nonconvex and discrete mixed integer program

Challenge: augmented Lagrangian methods typically do not converge

Part III: Distributed Computing Methods through ADMM

Augmented Lagrangian Relaxation

C Relaxing
$$\delta_{a,i,t} = \delta_{b,i,t}, \delta_{a,j,t} = \delta_{b,j,t}, \forall t \in T$$

Averaging
$$\overline{\delta}_{i,t} = \sum_{a \in \Phi(i)} \delta_{a,i,t} / |\Phi(i)|$$

Augmented model

$$\begin{split} \min_{\boldsymbol{\xi}_{a}^{f}} \ L_{a}\left(\boldsymbol{\xi}_{a}^{f}, \boldsymbol{\lambda}_{a}, \overline{\boldsymbol{\delta}}\right) & L_{a}\left(\boldsymbol{\xi}_{a}^{f}, \boldsymbol{\lambda}_{a}, \overline{\boldsymbol{\delta}}\right) = \sum_{i \in \mathbb{N}_{a}^{BB} \cup \mathbb{N}_{a}^{BB}} \sum_{t \in \mathbb{T}} \left[\lambda_{a,i,t}\left(\delta_{a,i,t} - \overline{\delta_{i,t}}\right) + \frac{\rho}{2}\left(\delta_{a,i,t} - \overline{\delta_{i,t}}\right)^{2} \right] \\ & + \max_{\tilde{\mathbf{P}}_{a}^{w} \in \mathbb{U}_{a}^{w}} \left[\min_{\boldsymbol{\xi}_{a}^{s} \in \Omega_{a}^{s}\left(\boldsymbol{\xi}_{a}^{f}, \tilde{\mathbf{P}}_{a}^{w}\right)} C_{a}^{ED}\left(\boldsymbol{\xi}_{a}^{s}, \tilde{\mathbf{P}}_{a}^{w}\right) \right] \\ & = C_{a}'\left(\boldsymbol{\xi}_{a}^{f}, \boldsymbol{\lambda}_{a}, \overline{\boldsymbol{\delta}}\right) + \max_{\tilde{\mathbf{P}}_{a}^{w} \in \mathbb{U}_{a}^{w}} \left[\min_{\boldsymbol{\xi}_{a}^{s} \in \Omega_{a}^{s}\left(\boldsymbol{\xi}_{a}^{f}, \tilde{\mathbf{P}}_{a}^{w}\right)} C_{a}^{ED}\left(\boldsymbol{\xi}_{a}^{s}, \tilde{\mathbf{P}}_{a}^{w}\right) \right] \end{split}$$

Overall Algorithm Scheme

- Using ADMM, each area can be computed independently
 - distributed computing and privacy protection
- For a single area problem: two-stage robust optimization
 - column and constraint generation method
 - finitely convergent for a polyhedron uncertainty set
- Integrated ADMM+CCG (IAC) Solution Method
 - Multi-area robust tie-line scheduling: ADMM+CCG converges to optimal value
 - Multi-area robust generation unit and tie-line scheduling: convergence is NOT guaranteed
 - Computational enhancements? Speed and convergence

Fast Computing

- Warm Start (WS):
 - select initial values of the first-stage variables and dual variables using the deterministic version
- Scenario Retaining (SR):
 - CCG is repeatedly called within ADMM framework
 - Keep and re-use existing scenarios generated in previous CCG calls
- Scenario Discard (SD):
 - Remove non-critical scenarios to maintain a small pool
 - Dynamically manage a scenario pool through a changing threshold
- SR and SD are key steps in distributed computation of Robust Optimization

Convergence Issue from UC

 Non-convergence due to non-convex structure from binary commitment decisions

- Alternating optimization procedure to ensure convergence (heuristically)
 - Alternatively computing with boundary phase angles or commitment status are fixed
 - A repeated commitment status indicates termination
 - Finitely converged



Part IV: Numerical Study and Demonstrations

Two-Area 12-Bus Interconnected System



Tie-line Flows



IAC Performance for Tie-line Scheduling



Computational Enhancement Strategies

	Strategy			Initial	nitialization ADMM							
Case		SR			Time				C&CG			Total
	WS		SD	# iter.		# iter.	# solution		Time (s)			time (s)
					(s)		MP	SP	MP	SP	SD	
M0	-	-	-	-	-	120	516	469	20.1	33.3	-	53.4
M1	-	\checkmark	-	-	-	120	151	152	223.2	11.0	-	234.2
M2	-	\checkmark	\checkmark	-	-	120	179	179	16.6	13.7	5.8	36.1
M3	\checkmark	-	-	21	5.1	41	82	82	0.9	4.3	-	10.3
M4	\checkmark	\checkmark	\checkmark	21	5.2	41	41	41	0.5	1.5	0.5	7.7

*M0: IAC without enhancement

Coordination Effect with Unit Commitment

slightly higher than the centralized solution by 0.11%

SCHEDULING RESULTS ON 12-BUS SYSTEM

	Proposed IAC					Centralized C&CG				Non-coordinated								
Hour	A	rea	1	Are	ea 2	Tie	A	rea	1	Are	ea 2	Tie	A	rea	1	Are	ea 2	Tie
	G1	G2	G3	G4	G5	(MW)	G1	G2	G3	G4	G5	(MW)	G1	G2	G3	G4	G5	(MW)
1	1	0	0	1	0	-10.4	1	0	0	1	0	-7.9	1	0	0	1	0	0.0
2	1	0	0	1	0	-1.9	1	0	0	1	0	16.1	1	0	0	1	0	0.0
3	1	0	0	1	0	5.4	1	0	0	1	0	18.1	1	0	0	1	0	0.0
4	1	0	0	1	0	11.8	1	0	0	1	0	24.4	1	0	0	1	0	0.0
5	1	0	0	1	0	11.6	1	0	0	1	0	25.3	1	0	0	1	0	0.0
6	1	0	0	1	0	12.6	1	0	0	1	0	21.4	1	0	0	1	0	0.0
7	1	0	0	1	0	25.9	1	0	0	1	0	25.9	1	0	0	1	1	0.0
8	1	0	0	1	0	16.7	1	0	0	1	0	16.5	1	0	0	1	1	0.0
9	1	0	0	1	0	-0.3	1	0	0	1	0	19.9	1	0	0	1	1	0.0
10	1	0	0	1	0	-3.7	1	0	0	1	0	12.7	1	0	0	1	1	0.0
11	1	0	1	1	1	3.0	1	0	1	1	0	11.0	1	0	1	1	1	0.0
12	1	0	1	1	1	2.0	1	0	1	1	0	7.6	1	0	1	1	1	0.0
13	1	0	1	1	1	2.7	1	0	1	1	0	6.7	1	0	0	1	1	0.0
14	1	0	1	1	1	3.4	1	0	1	1	0	7.6	1	0	0	1	1	0.0
15	1	0	1	1	1	1.7	1	0	1	1	0	6.1	1	0	1	1	1	0.0
16	1	0	0	1	0	-8.4	1	0	0	1	0	-0.6	1	0	0	1	0	0.0
17	1	0	0	1	0	-11.4	1	0	0	1	0	-8.4	1	0	0	1	0	0.0
18	1	0	0	1	0	-2.7	1	0	0	1	0	-2.0	1	0	0	1	0	0.0
19	1	0	0	1	0	-5.2	1	0	0	1	0	-9.9	1	0	0	1	0	0.0
20	1	0	0	1	0	-2.5	1	0	0	1	0	-34.3	1	0	0	1	0	0.0
21	1	0	0	1	0	-6.8	1	0	0	1	0	-9.0	1	0	0	1	0	0.0
22	1	0	0	1	0	9.3	1	0	0	1	0	4.6	1	0	0	1	0	0.0
23	1	0	0	1	0	7.1	1	0	0	1	0	0.4	1	0	0	1	0	0.0
24	1	0	0	1	0	-2.0	1	0	0	1	0	4.4	1	0	0	1	0	0.0
Obj.	\$150,419					\$150,243			\$153,645									

Performance in Large Systems

System	Areas	Units	Buses	Int. Lines	Tie-lines
2A-RTS	2	66	48	76	3
3A-RTS	3	99	73	115	5
118-Bus	3	79	118	174	12

IAC for Tie-line Scheduling

SIMULATION RESULTS ON LARGE-SCALE TEST SYSTEMS										
Unc	ertainty l	Budget	$\Gamma = 0$	Γ=6	Γ = 12	Γ = 24				
	Contr	Obj. (\$)	5,982,249	6,104,215	6,182,390	6,203,326				
	Cenu.	Time (s)	2.4	15.2	24.0	9.8				
2A PTS		Obj. (\$)	5,982,497	6,118,620	6,198,730	6,226,089				
ZA-KIS	IAC	# iter.	101	103	103	107				
	IAC	Time (s)	17.9	92.1	79.4	98.7				
		Error (%)	0.00	0.24	0.26	0.37				
	Canta	Obj. (\$)	9,151,653	9,333,234	9,459,759	9,510,397				
	Cenu.	Time (s)	3.6	33.6	106.7	9.9				
2A DTS	IAC	Obj. (\$)	9,151,694	9,348,872	9,468,978	9,512,639				
JA-KIS		# iter.	226	320	316	339				
		Time (s)	34.4	830.4	501.9	676.5				
		Error (%)	0.00	0.17	0.10	0.02				
	Casta	Obj. (\$)	2,249,048	2,250,762	2,252,142	2,255,044				
118-Bus	Centr.	Time (s)	4.4	4.6	4.7	4.2				
		Obj. (\$)	2,249,120	2,251,392	2,253,339	2,257,712				
	TAC	# iter.	491	535	573	600				
	IAC	Time (s)	197.1	285.1	325.2	348.3				
		Error (%)	0.00	0.03	0.05	0.12				

IAC for Generation Unit and Tie-line Scheduling

SIMULATION RESULTS ON LARGE-SCALE TEST SYSTEMS										
Uncer	rtainty B	udget	$\Gamma = 0$	$\Gamma = 6$	$\Gamma = 12$	$\Gamma = 24$				
	Centr	Obj. (\$)	1,014,164	1,018,237	1,039,095	1,044,114				
	Cenu.	Time (s)	597	3,005	1,617	1,451				
2A-RTS		Obj. (\$)	1,019,117	1,034,468	1,055,462	1,061,706				
	IAC	Time (s)	127	3,642	4,186	788				
		Gap (%)	0.49	1.59	1.58	1.66				
	Centr.	Obj. (\$)	1,529,977	1,535,347	1,565,097	1,576,184				
		Time (s)	25	4,196	5,876	4,821				
3A-RTS	IAC	Obj. (\$)	1,541,769	1,562,666	1,583,698	1,591,784				
		Time (s)	125	2,360	1,757	3,486				
3A-RTS		Gap (%)	0.77	1.78	1.19	0.99				
	Contr	Obj. (\$)	1,101,662	1,111,969	1,112,181	1,112,431				
118-Bus	Centr.	Time (s)	19	922	316	867				
	IAC	Obj. (\$)	1,111,116	1,113,739	1,117,006	1,119,490				
		Time (s)	172	4,775	918	437				
3A-RTS		Gap (%)	0.85	1.06	1.34	1.54				

Part V: Conclusions

Observations and Conclusions

- An integrated ADMM+CCG computing method
 - Supporting information and privacy protection in handling uncertainties
 - Advanced enhancement strategies for fast computation
 - New strategies to ensure convergence in non-convex structures
- Coordination plays a critical role in multi-area grid performance
 - For tie-line scheduling, IAC performs (almost) the same as centralized method
 - For commitment and tie-line scheduling, IAC significantly outperforms non-coordinated control
- Future Improvement
 - Economic implications from IAC computation
 - Novel algorithmic improvement to support fast computing