

Power-Capacity and Ramp-Capability Reserves for Wind Integration in Power-based Unit Commitment¹

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FERC: Increasing Market and Planning Efficiency
through Improved Software

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¹G. Morales-España, R. Baldick, J. García-González, and A. Ramos, "Power-Capacity and Ramp-Capability Reserves for Wind Integration in Power-Based UC," *IEEE Transactions on Sustainable Energy*, vol. 7, no. 2, pp. 614–624, Apr. 2016

Outline

- 1 Introduction
- 2 Power-Capacity and Ramp-Capability Reserves
 - Why Ramp-Capability Reserves?
 - Reserves Logic
- 3 Case Studies
- 4 Conclusions

Reserves in UC

- Wind & Solar introduce uncertainty \Rightarrow more difficult planning

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 - Implicit reserves
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- Optimal quantity of reserves must be scheduled
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- **Stochastic UCs**
 - Implicit reserves
 - \uparrow computational burden
- **Reserve-based Deterministic UC**
 - Explicit reserves
 - \downarrow computational burden

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1 Introduction

2 Power-Capacity and Ramp-Capability Reserves

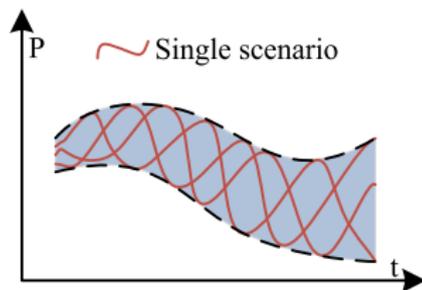
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Why Ramp-Capability Reserves?

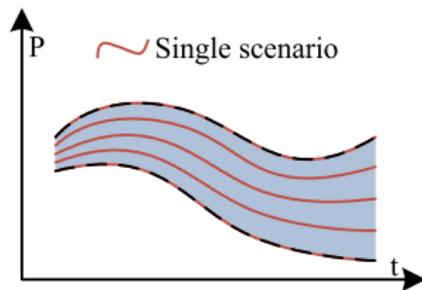
Stochastic



Implicitly guarantees availability of resources for a capacity and ramp range

Why Ramp-Capability Reserves?

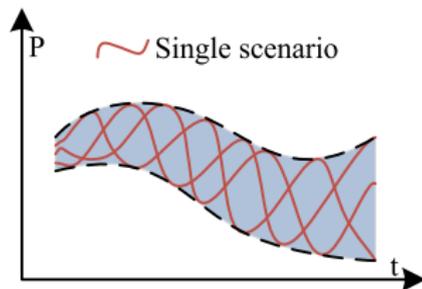
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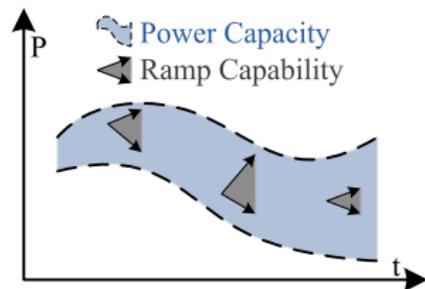
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Capacity & Ramp Reserves

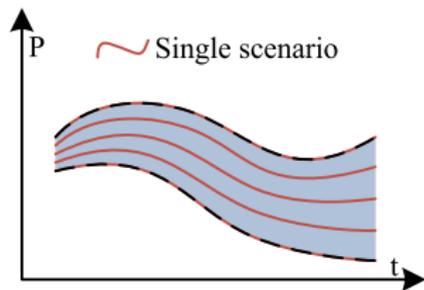


Explicitly guarantee availability of resources for capacity and ramp reserves requirements

Need for a clear difference between
Power-Capacity and **Ramp-Capability** Requirements

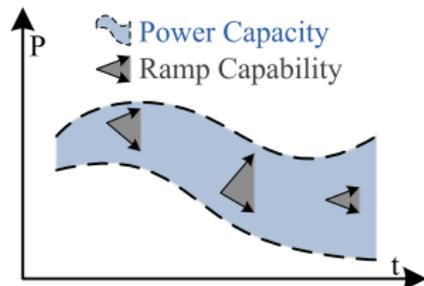
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 - **Adds robustness with this worst-case scenario**

Single-level MIP for the Robust UC

- By considering dispatchable wind, the Robust UC becomes²

$$\begin{array}{ll}
 \min & \mathbf{b}^\top \mathbf{x} + \max \min \mathbf{c}^\top \mathbf{p} \\
 \text{s.t.} & \mathbf{F}\mathbf{x} \leq \mathbf{f}, \mathbf{x} \text{ is binary} \\
 & \mathbf{H}\mathbf{p} + \mathbf{J}\mathbf{w} \leq \mathbf{h}, \forall \xi \in \Xi \\
 & \mathbf{A}\mathbf{x} + \mathbf{B}\mathbf{p} \leq \mathbf{g}, \forall \xi \in \Xi \\
 & \mathbf{w} \leq \xi, \forall \xi \in \Xi
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where uncertainty set Ξ is defined by $\xi_{bt} \in [\underline{w}_{bt}, \overline{w}_{bt}] \forall t \in \mathcal{T}, b \in \mathcal{B}^w$

- Which is a considerably simpler problem, **we avoid**
 - The local optimum of the bilinear program
 - Further complexity when trying to solve the **two-level bilinear** + **MIP**

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- Which is a considerably simpler problem, we avoid
 - The local optimum of the bilinear program
 - Further complexity when trying to solve the two-level bilinear + MIP
- The worst-case wind scenario can be known a priori
 - this key worst-case scenario gives robustness to the UC solution

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 - Adds robustness with this worst-case scenario
 - **readjust power-capacity reserves requirement**
- The procured **ramp-capability** \leq **power-capacity reserves**
 - \Rightarrow **readjust ramp-capability reserves requirement**

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 - 54 thermal units, 3 wind farms, 186 transmission lines
 - 24 hours time span
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 - DetRes³: UC with traditional power-capacity reserves
 - Stch: Stochastic UC
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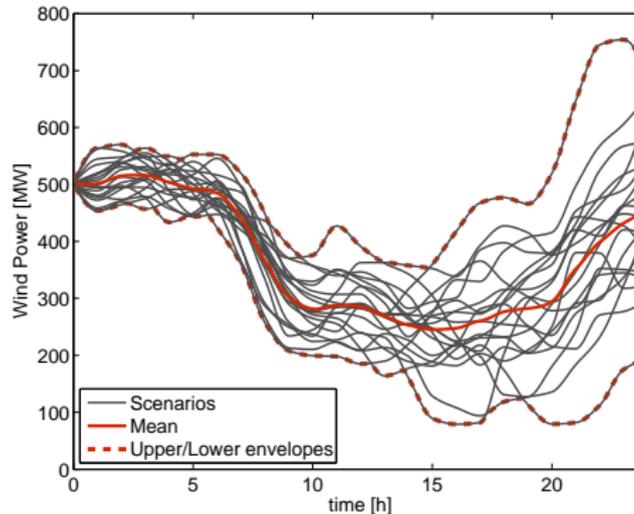
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- 3 Power-based UC formulations implemented:
 - DetRes³: UC with traditional power-capacity reserves
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 - ResRPC⁴: UC with Power-capacity & Ramp-capability reserves
- All problems solved with Cplex 12.6.0, stop criteria:
 - 0.05% opt. tolerance or 2h time limit

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Scheduling & Evaluation Stages

- **Scheduling Stage: 20 in-sample scenarios**
 - Obtains **hourly** commitment decisions for all units
 - by solving **hourly network-constrained UCs**
 - Reserves **obtained from** the in-sample scenarios:



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- **Evaluation Stage: 200 out-of-sample scenarios**
 - **5 min** dispatch decisions for all units
 - by solving **5-min network-constrained optimal dispatch**
 - Penalizations:
 - Demand-balance violation costs: 10000 \$/MWh
 - Network violation costs: 5000 \$/MWh

Scheduling Performance

Scheduling (hourly)		
	UC Costs [†] [k\$]	# SU
DetRes	55.49	16
Stch	54.77	12
ResRPC		

[†] Commitment cost

- Stch optimized the level of reserves by dispatching wind
 - Scheduled less units
 - Lower UC costs than DetRes

Scheduling Performance

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 - Scheduled less units
 - Lower UC costs than DetRes
- ResRPC also optimized the level of reserves

Evaluation: 200 out-of-sample scenarios

	Scheduling (hourly)		Real-time dispatch (5-min)	
	UC Costs [†] [k\$]	# SU	Costs* [k\$]	# Tot. Viol.
DetRes	55.49	16	857.20	611
Stch	54.77	12	808.97	259
ResRPC	51.98	14		

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* average dispatch + penalty costs of the 200 scenarios

- Compared with **Stch**, **DetRes** presented
 - **6%** higher average dispatch costs
 - **2.4x** more violations

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 - **5%** lower average dispatch costs
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 - **mainly due to the robustness of ResRPC**

Evaluation: 20 in-sample scenarios

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	UC Costs [†] [k\$]	# SU	Costs* [k\$]	# Tot. Viol.
DetRes	55.49	16	803.46	611
Stch	54.77	12	768.79	12
ResRPC	51.98	14	770.86	1

[†] Commitment cost

* average dispatch + penalty costs of the 20 scenarios

- Compared with ResRPC, Stch presented
 - **0.3%** lower average dispatch costs
 - even with **12x** more violations

Power- vs. Energy-based UC: o-of-s evaluation

	Scheduling (hourly)		Real-time dispatch (5-min)	
	UC Costs [†] [k\$]	# SU	Costs* [k\$]	# Tot. Viol.
DetRes	55.49	16	857.2	611
Stch	54.77	12	808.97	259
ResRPC	51.98	14	770.82	2
E-Stch	33.73	10	1166.67	1382

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- Compared with a traditional Energy-based stochastic UC **E-Stch**, **DetRes**
 - presented **27%** lower average dispatch costs
 - presented **56%** less violations

Power- vs. Energy-based UC: o-of-s evaluation

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- Compared with a traditional Energy-based stochastic UC **E-Stch**, **DetRes**
 - presented **27%** lower average dispatch costs
 - presented **56%** less violations
 - and solved **23.6x** faster

Computational Burden

- The power-based formulations [DetRes](#)⁵, [Stch](#), [ResRPC](#)⁶
 - include startup and shutdown power trajectories⁷
 - are built upon a [convex-hull](#)⁸

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Computational Burden

- The power-based formulations [DetRes](#)⁵, [Stch](#), [ResRPC](#)⁶
 - include startup and shutdown power trajectories⁷
 - are built upon a [convex-hull](#)⁸
- Two energy-based UC formulations:
 - [E-Stch](#)⁹, also based on a [convex-hull](#)¹⁰
 - [TE-Stch](#)¹¹, Traditional formulation, var. startup costs¹²

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Computational Burden

	MIP Time [s]	LP relaxation [s]
DetRes	8.8	0.34
Stch	867.8	38.1
ResRPC	90.4	16.8
E-Stch	206.5	22.1
TE-Stch		

- ResRPC solved
 - **10.3x** slower than DetRes
 - but **9.6x** faster than Stch
 - and **2.3x** faster than E-Stch

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- E-Stch could not reach 0.05% within the **2h time limit**

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 - and **2.3x** faster than E-Stch
- E-Stch could not reach 0.05% within the **2h time limit**
- ResRPC, Stch and DetRes TE-Stch solved the MIP before E-Stch solved the LP

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 - better represents requirements to accommodate wind
 - outperforms a power-based UC with traditional reserves
 - can beat a stochastic power-based UC
 - is deterministic \Rightarrow \downarrow computational burden
- Compared with traditional energy-based UC, the deterministic power-based UCs
 - presented lower average costs and fewer violations in the out-of-sample 5/min dispatch evaluation
 - while solving MIP problems faster

Questions

Thank you for your attention

Contact Information:
g.a.moralesespama@tudelft.nl

For Further Reading



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